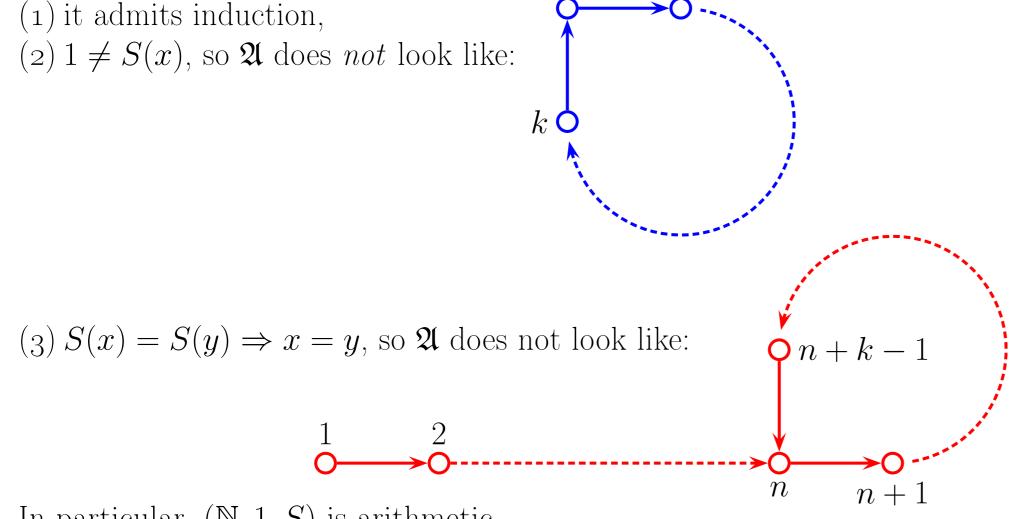
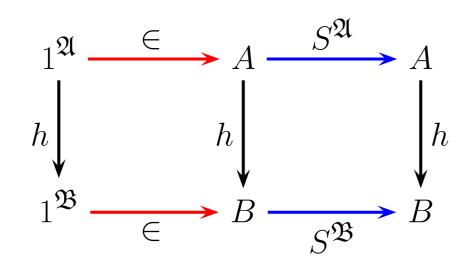
Numbers and sets David Pierce May, 2010 • () • ) In  $\{a, b, -, +, \cdot\}$ , the term  $+ \cdot + a b a - + ab$ , or more conventionally By an **iterative algebra**, I mean an ordered triple (A, 1, S), or  $\mathfrak{A}$ , where (1) A is a set;  $(2) \ 1 \in A;$ (3)  $S: A \to A$ .  $((a+b)\cdot a) + -(a+b).$ **Example.**  $A = \mathbb{N} = \{1, 2, 3, ...\}$  and S(n) = n + 1. An iterative algebra  $\mathfrak{A}$  can be conceived as a **directed graph**, where (1) A is the set of nodes, (2) 1 is a particular node, and (3) each pair (x, S(x)) is an **arrow** from x to S(x).  $\begin{array}{c|c|c|c|c|c|} & \text{and} & \\ + & & + \\ \end{array}$ corresponds to either of . Now let The iterative algebra  $\mathfrak{A}$  admits induction if it has no proper subalgebra, so it does *not* look like:

Then **A** is **arithmetic** if



In particular,  $(\mathbb{N}, 1, S)$  is arithmetic.

In general,  $\mathfrak{A}$  admits recursion if, for every iterative algebra  $\mathfrak{B}$ , there is a unique homomorphism h from  $\mathfrak{A}$  to  $\mathfrak{B}$ :



**Theorem** (Dedekind). An iterative algebra admits recursion if and only if it is arithmetic; in particular, all such iterative algebras are isomorphic. Also  $\mathbb{N}$  is well-ordered by the relation <, defined recursively by

$$x \not< 1,$$
  $x < n+1 \Leftrightarrow x \leqslant n.$ 

Functions can be defined on  $\mathbb{N}$  by well-ordered recursion: the simplest example is h, given by

$$\begin{vmatrix} + \\ \cdot \\ \cdot \\ b \end{vmatrix} = (\{a\}, \{b\}, +), \qquad \begin{vmatrix} + \\ + \\ \cdot \\ a \end{pmatrix} = \left( \left\{ a, b, (\{a\}, \{b\}, +) \right\}, \{a\}, \cdot \right),$$

and so on; then (\*) is as depicted below. If F is n-ary in  $\mathscr{S}$  as before,  $X_k$  is a set when k < n, and  $X = (X_0, \ldots, X_{n-1}, F)$ , define

$$\operatorname{pred}_k(X) = X_k,$$
  $\operatorname{pred}(X) = X_0 \cup \cdots \cup X_{n-1},$   
 $Y \in X \Leftrightarrow Y \in \operatorname{pred}(X)$ 

(here 'pred' is for *predecessor*); say that X is *k*-transitive if

 $Y \in \operatorname{pred}_k(X) \Rightarrow \operatorname{pred}(Y) \subseteq \operatorname{pred}_k(X).$ 

Let  $\mathbf{ON}_{\mathscr{S}}$  comprise those X such that

 $= (\{a\}, \{b\}, +),$ 

a

(1)  $\overline{X} = (\overline{X}_0, \dots, \overline{X}_{n-1}, F)$  for some *n*-ary *F* in  $\mathscr{S}$ , for some *n* in  $\boldsymbol{\omega}$ ; and each  $X_k$  is nonempty;

(2) each element Y of  $\operatorname{pred}(X)$  is  $(Y_0, \ldots, Y_{m-1}, G)$  for some *m*-ary G in  $\mathscr{S}$ , for some m in  $\omega$ ; and each  $Y_{\ell}$  is nonempty;

(3) X is k-transitive for each k;

(4) each element of  $\operatorname{pred}(X)$  is  $\ell$ -transitive for each  $\ell$ ;

 $(5) \in'$  directs each set  $\operatorname{pred}_k(X)$  (finite subsets have upper bounds);

(6)  $\in'$  directs each set  $\operatorname{pred}_{\ell}(Y)$  for each Y in  $\operatorname{pred}(X)$ ;

 $(7) \in '$  is well-founded on pred(X) (nonempty subsets have minimal

elements).

(\*)

Call an element X of  $\mathbf{ON}_{\mathscr{S}}$  a **limit** if some  $\operatorname{pred}_k(X)$  has no maximal element with respect to  $\in'$ . Let  $\omega_{\mathscr{S}}$  consist of those X in  $ON_{\mathscr{S}}$  such that neither X nor any element of  $\operatorname{pred}(X)$  is a limit.

**Theorem.** The relation  $\in'$  is well-founded on  $ON_{\mathscr{S}}$ , and

 $X \in \mathbf{ON}_{\mathscr{S}} \Rightarrow \operatorname{pred}(X) \subseteq \mathbf{ON}_{\mathscr{S}}.$ 

 $h(n) = \{h(x) \colon x < n\}.$ 

Then h is a bijection from N onto  $\omega$ , the set of **von Neumann** natural numbers. The first five of these -0, 1, 2, 3, and 4 are illustrated above.

The class **ON** of von Neumann **ordinals** comprises each set that

(1) is well-ordered by membership  $(\in)$ ,

(2) is **transitive** (its members are also subsets).

Then **ON** itself is well-ordered by membership and is transitive; it is an iterative algebra with respect to  $\emptyset$  and  $x \mapsto x \cup \{x\}$ ; and  $\omega$  is a subalgebra of **ON** and is a **free algebra**.

All of the foregoing generalizes to an arbitrary algebraic signature,  $\mathscr{S}$ . One free algebra in  $\mathscr{S}$  is the **term algebra:** the smallest set of strings that, for each n in  $\omega$ , for each n-ary symbol F in  $\mathscr{S}$ , is closed under the concatenation

$$(t_0,\ldots,t_{n-1})\longmapsto Ft_0\cdots t_{n-1}.$$

Terms can be written as **labelled trees** or (refinements of) **Hasse diagrams**:

The class  $\mathbf{ON}_{\mathscr{S}}$  is an  $\mathscr{S}$ -algebra with respect to the operations

 $F(X_0, \ldots, X_{n-1}) = (\operatorname{pred}(X_0) \cup \{X_0\}, \ldots, \operatorname{pred}(X_{n-1}) \cup \{X_{n-1}\}, F);$ 

and  $\omega_{\mathscr{S}}$  is a free subalgebra.

