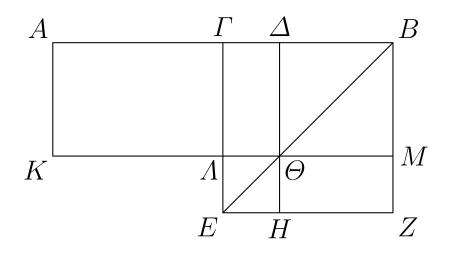
D. Pierce, Vector-spaces over unspecified fields, Nijmegen, 2006From Euclid of Alexandria, *The Elements*, Proposition II.5:

Εὐθεῖα γάρ τις ἡ ΑΒ τετμήσθω εἰς μὲν ἴσα κατὰ τὸ Γ, εἰς δὲ ἀνισα κατὰ τὸ Δ[·] λέγω, ὅτι τὸ ὑπὸ τῶν ΑΔ, ΔΒ περιεχόμενον ὀρθογώνιον μετὰ τοῦ ἀπὸ τῆς ΓΔ τετραγώνου ἴσον ἐστι τῷ ἀπὸ τῆς ΓΒ τετραγώνῳ.



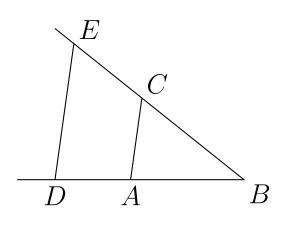
'For, let a straight [line] AB be cut into equal [segments] at Γ , and unequal at Δ ; I say that the rectangle bounded by $A\Delta$, ΔB , with the square on $\Gamma\Delta$, is equal to the square on ΓB .'

We might say

$$(x+y)(x-y) + y^2 = x^2,$$

where x, y are the lengths of $A\Gamma$, $\Gamma\Delta$ respectively; maybe Euclid too (Descartes, Rules for the Direction of the Mind, 4).

From René Descartes, second page of *The Geometry*:



«Soit par exemple AB l'vnité, & qu'il faille multiplier BD par BC, ie n'ay qu'a ioindre les poins A & C, puis tirer DE parallele a CA, & BE est le produit de cete Multiplication.»

'For example, suppose AB is unity, and that one must multiply BD by BC; I need only join points A and C, then draw DE parallel to CA, and BE is the product of this multiplication.'

Take B as an origin of **vectors.** Then D is the multiple of A by a scalar, say [D:A]. Because $E - D \parallel C - A$, we have

$$[D:A] * C = E$$
 & $[D:A] = [E:C].$

Precisely, a **vector-space** is a two-sorted structure (V, K, *), where:

V is an abelian group; K is a field; $*: K \times V \to V$;

 $(t \mapsto (\boldsymbol{x} \mapsto (t * \boldsymbol{x}))) \in \operatorname{Hom}(K, (\operatorname{End}(V), \circ)).$

So the signature of (V, K, *) is $\{+, -, 0\} \amalg \{+, -, \cdot, 0, 1\} \amalg \{*\}$; there are variables \boldsymbol{x} for vectors, and t for scalars.

Abraham Robinson (*Complete Theories*, Amsterdam, 1956) showed model-completeness of the theory of non-trivial vector-spaces over a *fixed* scalar field: 'partial' model-completeness.

If $n \in \{1, 2, 3, ..., \infty\}$, let T_n be the theory of *n*-dimensional vectorspaces. Andrey Kuzichev (*Z. Math. Logik Grundlag. Math.*, 1992) showed elimination of quantified *vector*-variables in T_n . Hence, if *U* is a complete field-theory, then $T_n \cup U$ is complete.

 T_n has $\exists \forall$ axioms for raising dimension. 'Hence:'

 $T_n \cup U$ is not inductive $(\forall \exists)$; much less model-complete (unless n = 1). We can remedy this 'problem' with a predicate for parallelism. The vector-space (V, K, *) expands to $(V, K, *, \parallel)$, where

$$\boldsymbol{x} \parallel \boldsymbol{y} \iff \exists s \; \exists t \; (s \ast \boldsymbol{x} + t \ast \boldsymbol{y} = \boldsymbol{0} \land (s \neq 0 \lor t \neq 0)).$$

If $\dim_K V > 1$, then $(V, K, *, \parallel)$ is interpretable uniformly in the reduct (V, \parallel) , and

$$(V,K,*,\parallel)\subseteq (W,L,*,\parallel)\iff (V,\parallel)\subseteq (W,\parallel).$$

(By contrast, three-sorted structures (G, N, G/N), where $N \triangleleft G$, are interpretable in the reducts (G, N); but

$$(\mathbb{Z}, 4\mathbb{Z}, \mathbb{Z}_4) \nsubseteq (\mathbb{Z}, 2\mathbb{Z}, \mathbb{Z}_2) \& (\mathbb{Z}, 4\mathbb{Z}) \subseteq (\mathbb{Z}, 2\mathbb{Z}).)$$

The theory VS₂ of the expansions $(V, K, *, \parallel)$ is inductive.

Hence the theory of the reducts (V, \parallel) is inductive.

The model-companion of VS_2 is the theory VS_2^* of **two-dimensional** vector-spaces over **algebraically closed** fields. In particular, VS_2^* is model-complete:

Replace the binary \parallel with an *n*-ary predicate \parallel^n for linear dependence; in the new signature, let VS_n be the theory of vector-spaces.

Theorem. The existentially closed models of VS_n are n-dimensional (over algebraically closed fields; since these models compose an elementary class, their theory is the model-companion of VS_n).

Proof. Say $K \subset L$ (both fields), and $\dim_K L \ge n+1$.

I say that $(K^{n+1}, K, \|^n)$ embeds in $(L^n, L, \|^n)$:

Suppose $\{a_0, \ldots, a_{n-1}, 1\}$ from L is linearly independent over K. Let

$$A = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ & \ddots & \\ 0 & 0 & 1 \\ -a_0 & -a_1 & \cdots & -a_{n-1} \end{pmatrix} = \left(\frac{I_n}{-a}\right).$$

The embedding is $\boldsymbol{x} \mapsto \boldsymbol{x} \cdot A$, taking the rows of I_{n+1} to the rows of A:

Indeed, write n elements of K^{n+1} as the rows of $(U | \boldsymbol{v})$. These rows are dependent if and only if

$$0 = \det\left(\frac{U \mid \boldsymbol{v}}{\boldsymbol{a} \mid 1}\right) = \det(U - \boldsymbol{v} \cdot \boldsymbol{a}).$$

But $U - \boldsymbol{v} \cdot \boldsymbol{a} = (U | \boldsymbol{v}) \cdot \left(\frac{I_n}{-\boldsymbol{a}}\right)$, whose rows are the images in L^n of the rows of $(U | \boldsymbol{v})$.

Compare with structures (L, K, P^n) , where P^n is *n*-ary algebraic dependence. From a standard counterexample:

$$(K(a_0,\ldots,a_n),K,P^n) \subset (K(a_0,\ldots,a_n,b,c),K(b,c),P^n),$$

where $\dim(a_0 \cdots a_n/K) = n + 1$, and (b, c) is a generic solution to $\sum_{i=0}^n a_i x^i = y$, so that $\dim(a_0 \cdots a_n/Kbc) = n$.

